

Event B: Sets, Relations, Functions, Arithmetic¹

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¹Many slides borrowed from J. R. Abrial: see http://wiki.event-b.org/index.php/Event-B_Language

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Set theory: membership



Set theory: membership



- Event-B formal reasoning is built based on:
 - First-order logic inference rules (seen).
 - Set theory (to be touched upon).
- Set theory as a foundation for relations, functions (and, therefore, data structures).
 - Proofs often reduced to proving goals on sets.
- We will briefly see how this is intuitively done.

- A set is a well-defined collection of distinct objects.
- Set theory is primary concerned the membership predicate

$\mathsf{E}\in\mathsf{S}$

• *E* is an expression, *S* is a set.

Set theory: basic constructs



Set theory: basic constructs Definitions



There are three basic constructs in set theory:

Cartesian product	S imes T
Power set	$\mathbb{P}(S)$
Comprehension 1	$\{x\cdot x\in S\ \wedge\ P(x)\mid F(x)\}$
Comprehension 2	$\set{x\mid x\in S\ \land\ P(x)}$

S and T are sets, *x* is a variable, P is a predicate, F is an expression.

 $\{1,2,3\} \times \{a,b\} = \{1 \mapsto a, 1 \mapsto b, 2 \mapsto a, 2 \mapsto b, 3 \mapsto a, 3 \mapsto b\}$

 $\mathbb{P}(\{1,2,3\}) = \{\{1,2,3\},\{1,2\},\{1,3\},\{2,3\},\{1\},\{2\},\{3\},\emptyset\}$

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Defined by equivalences

 $E \mapsto F \in S \times T \equiv E \in S \wedge F \in T$ $S \in \mathbb{P}(T) \equiv \forall x \cdot x \in S \Rightarrow x \in T$ $E \in \{x \mid x \in S \wedge P(x)\} \equiv E \in S \wedge P(E)$ $E \in \{x \cdot x \in S \wedge P(x) \mid F(x)\} \equiv \exists x \cdot x \in S \wedge P(x) \wedge E = F(x)$

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Set theory: basic constructs Examples



Operations on sets

$S \subseteq T$	\equiv	$S \in \mathbb{P}(T)$
S = T	\equiv	$S \subseteq T \land T \subseteq S$
$S \cup T$	\equiv	$\{x \mid x \in S \lor x \in T\}$
$S\cap T$	\equiv	$\{x \mid x \in S \land x \in T\}$
$S \setminus T$	\equiv	$\{x \mid x \in S \land x \notin T\}$
$E \in \{a, \ldots, z\}$	\equiv	$E = a \lor \ldots \lor E = z$
$E \in \emptyset$	\equiv	\perp

- software Polifecnic
- Operators based on membership and logic operations.
- Note: $E \notin T \equiv \neg (E \in T)$.
- Also: generalized / conditional union and intersection (see reference cards).

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Reminder: A \mapsto B is a tuple.
It is sometimes written as (A, B) in other formalisms.
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Shortcut: $m..n \equiv \{x \in \mathbb{Z} \mid m \le x \land x \le n\}$

 $\{x \mid x \in \{2,3,4,5\} \land x \mod 2 = 0\} = \{2,4\}$

 $\{x \cdot x \in \{2,3,4,5\} \land x \mod 2 = 1 \mid x^2\} = \{25,9\}$

• { $x \mid x \in \mathbb{N} \land x < 2$ } × 8..10 • { $n \cdot n \in \mathbb{N} \mid (0..n) \mapsto n$ } • { $x \cdot x \in 3..5 \mid x \mapsto x * x$ }

Binay relations



Types of relations



- A binary relation $r \in S \leftrightarrow T$ is a subset of their Cartesian product: $r \subseteq S \times T$
- Different syntax to highlight structure.
- $S \leftrightarrow T$: all (= the set of) the possible relations between S and T.
 - r would be one of them.

 $x \in dom(r) \equiv \exists y \cdot x \mapsto y \in r$ $y \in ran(r) \equiv \exists x \cdot x \mapsto y \in r$

• $r = \{1 \mapsto 10, 2 \mapsto 7, 2 \mapsto 11\}$

• $r \in 1..3 \leftrightarrow 7..11$

• $4 \mapsto 10 \notin r$

- $r \in \{\text{meat}, \text{fish}, \text{pasta}, \text{bacon}\} \leftrightarrow \{\text{carbs}, \text{protein}, \text{fat}\} \text{write a couple of relations.}$
- dom(r), ran(r), relation with S and T
- How many different *r* may there be?

 $r^{-1} \equiv \{ y \mapsto x \mid x \mapsto y \in r \}$

 $S \Leftrightarrow T \quad r \in S \Leftrightarrow T \land dom(r) = S$ Total Surjective $S \leftrightarrow T$ $r \in S \leftrightarrow T \land ran(r) = T$ $S \Leftrightarrow T$ $r \in S \Leftrightarrow T \land r \in S \Leftrightarrow T$ Both

Hint: sets and relations are very useful modeling tools!

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Operations on relations

Domain restriction	<i>S</i> ⊲ <i>r</i>	$\{x \mapsto y \in r \mid x \in S\}$
Domain subtraction	$S \lhd r$	$\{x \mapsto y \in r \mid x \notin S\}$
Range restriction	$r \rhd T$	$\{x\mapsto y\in r\mid y\in T\}$
Range subtraction	$r \triangleright T$	$\{x\mapsto y\in r\mid y\notin T\}$

Assume $Prey \in Animal \leftrightarrow Animal$.

We mean hunter \mapsto hunted. The syntax of the relation does not reveal its intended semantics.

- - $Mammal \triangleleft Prev$
 - Mammal *⊲* Prev
 - $Prey \triangleright Spiders$
 - Fish \triangleleft (Prey \triangleright Spiders)
 - Spiders \triangleleft (Prey \triangleright Spiders)



Operations on relations

Image	r[S]	$\{y \mid x \mapsto y \in r \land x \in S\}$
Composition	p; q	$\{x \mapsto z \mid x \mapsto y \in p \land y \mapsto z \in q\}$
Overriding	$p \Leftrightarrow q$	$q \cup (\mathit{dom}(q) \lhd p)$
Identity	id(S)	$\{x\mapsto x\mid x\in S\}$

Overriding:

- Take *q*, and add the tuples from *p* whose lhs are not already in *q*.
- Or, take *p* and add *q*, overriding the tuples with the same lhs.

Some useful results, definitions



$(r^{-1})^{-1}$	=	r
$\operatorname{dom}(r^{-1})$	=	ran(r)
$(S \lhd r)^{-1}$	=	$r^{-1} \triangleright S$
$(p;q)^{-1}$	=	$q^{-1}; p^{-1}$
p;(q;r)	=	(p;q);r
p ; $(q \cup r)$	=	$(p;q) \cup (p;r)$
(p; q)[S]	=	q[p[S]]
$r[S \cup T]$	=	$r[S] \cup r[T]$

 $r = r^{-1}$ symmetric $r \cap r^{-1} = \varnothing$ asymmetric $id(S) \subseteq r$ reflexive $r; r \subseteq r$ transitive

Functions

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Functions: one type of relations. Notation: $f(x) = y \equiv x \mapsto y \in f$.	Total function $(dom(f) = S)$ Partial function	$S \rightarrow T$ $S \rightarrow T$	
Every element in domain relates only to one element in range.	Injection if $f(x) = f(x)$ then		
to one element in range.	Injection: if $f(x) = f(y)$, then	x = y.	
$x \mapsto y \in f \land x \mapsto z \in f \Rightarrow y = z$	Partial injection	$S \rightarrowtail T$	
	Total injection	$S \rightarrowtail T$	
WD conditions:			
• $f \in S \rightarrow T$ • $x \in dom(f)$	Surjection: $f \in S \leftrightarrow T$, $ran(f)$	= T.	
	Partial surjection	S +⇒ T	
Using right type of function allows different proofs.	Total surjection	$S \twoheadrightarrow T$	
	Bijection	$S \rightarrowtail T$	

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An example of functions and relations: a strict society

Set-theoretic notation more readable than predicate calculus

 $r = r^{-1} \equiv \forall x, y \cdot x \in S \land y \in S \Rightarrow (x \mapsto y \in r \Leftrightarrow y \mapsto x \in r)$



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An example of functions and relations: a strict societ	y
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Every person is man or woman

 $men \subseteq PERSON$

- Every person is either a man or a woman.
- No person is man and woman at the same time.
- Only women have husbands, who must be a man.
- Woman have at most one husband.
- Men have at most one wife.
- Mother are married women.



Every person is man or woman No person is man and woman

 $men \subseteq PERSON$ women = $PERSON \setminus men$

An example of functions and relations: a strict society



Every person is man or woman No person is man and woman Women have husbands (men) At most one husband per woman Men at most one wife

 $men \subseteq PERSON$ *women* = *PERSON* \setminus *men*

 $husband \in women \rightarrowtail men$

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An example of functions and relations: a strict society

Every person is man or woman
No person is man and woman
Women have husbands (men)
At most one husband per woman
Men at most one wife
Mother are married women

$men \subseteq PERSON$	

women = $PERSON \setminus men$

 $husband \in women
arrow men$

mother \in *PERSON* \rightarrow dom(*husband*)

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An example of functions and relations: a strict society



Every person is man or woman	$men \subseteq PERSON$
No person is man and woman	women = $PERSON \setminus men$
Women have husbands (men)	
At most one husband per woman	$\mathit{husband} \in \mathit{women} \rightarrowtail \mathit{men}$
Men at most one wife	
Mother are married women	$mother \in PERSON \leftrightarrow dom(husband)$

Let us derive some relations (Double check with Rodin)

wife =	daughter =
spouse =	sibling =
<i>father</i> =	brother =
children =	

An example of functions and relations: a strict society



An example of functions and relations: a strict society



Every person is man or woman	$men \subseteq PERSON$
No person is man and woman	women = $PERSON \setminus men$
Women have husbands (men)	
At most one husband per woman	$\mathit{husband} \in \mathit{women} \rightarrowtail \mathit{men}$
Men at most one wife	
Mother are married women	$mother \in PERSON \leftrightarrow dom(husband)$

Let us derive some relations (Double check with Rodin)

wife = $husband^{-1}$ spouse = father = children =

daughter = sibling = *brother* =

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An example of functions and relations: a strict society

Every perso	n is man or woman	$men \subseteq PERSON$	
No person	is man and woman	women = $PERSON \setminus men$	
Women hav	/e husbands (men)		
At most one	e husband per woman	$\mathit{husband} \in \mathit{women} \rightarrowtail \mathit{men}$	
Men at mos	st one wife		
Mother are	married women	$mother \in PERSON \rightarrow dom(husband)$	1)

Let us derive some relations (Double check with Rodin)

wife = $husband^{-1}$
$spouse = husband \cup wife$
father = mother; husband
<i>children</i> =

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sibling =	
brother =	

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Every person is man or woman	$men \subseteq PERSON$
No person is man and woman	women = $PERSON \setminus men$
Women have husbands (men)	
At most one husband per woman	$\mathit{husband} \in \mathit{women} \rightarrowtail \mathit{men}$
Men at most one wife	
Mother are married women	$\textit{mother} \in \textit{PERSON} \rightarrow \textit{dom}(\textit{husband})$

Let us derive some relations (Double check with Rodin)

wife = $husband^{-1}$ $spouse = husband \cup wife$ father = children =

daughter = sibling = brother =

 \subseteq PERSON

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An example of functions and relations: a strict society	
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Every person is man or woman	$men \subseteq PERSON$
No person is man and woman	women = $PERSON \setminus men$
Women have husbands (men)	
At most one husband per woman	$\mathit{husband} \in \mathit{women} \rightarrowtail \mathit{men}$
Men at most one wife	
Mother are married women	$mother \in PERSON \leftrightarrow dom(husband)$

Let us derive some relations (Double check with Rodin)

wife = $husband^{-1}$ $spouse = husband \cup wife$ *father* = *mother*; *husband* children = $(mother \cup father)^{-1}$

daughter = sibling = brother =

An example of functions and relations: a strict society



An example of functions and relations: a strict society



Every person is man or woman No person is man and woman Women have husbands (men) At most one husband per woman Men at most one wife

 $men \subseteq PERSON$ women = $PERSON \setminus men$

Mother are married women

husband ∈ *women* →→ *men*

mother \in *PERSON* \rightarrow dom(*husband*)

Let us derive some relations (Double check with Rodin)

wife = $husband^{-1}$ $spouse = husband \cup wife$ father = mother: husband children = $(mother \cup father)^{-1}$ $daughter = children \lhd women$ sibling = brother =

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An example of functions and relations: a strict society

Every person is man or woman	$men \subseteq PERSON$
No person is man and woman	women = $PERSON \setminus men$
Women have husbands (men)	
At most one husband per woman	$\mathit{husband} \in \mathit{women} \rightarrowtail \mathit{men}$
Men at most one wife	
Mother are married women	$mother \in PERSON \leftrightarrow dom(husb)$

Let us derive some relations (Double check with Rodin)

wife = $husband^{-1}$ $spouse = husband \cup wife$ *father* = *mother*; *husband* children = $(mother \cup father)^{-1}$ $daughter = children \lhd women$ sibling = $(children^{-1}; children) \setminus id(PERSON)$ $brother = sibling \triangleright men$

Every person is man or woman No person is man and woman Women have husbands (men) At most one husband per woman Men at most one wife Mother are married women

Let us derive some relations (Double check with Rodin)

wife = $husband^{-1}$ $spouse = husband \cup wife$ father = mother; husband children = $(mother \cup father)^{-1}$

Properties

 $daughter = children \lhd women$ sibling = $(children^{-1}; children) \setminus id(PERSON)$ brother =

men \subset *PERSON*

women = $PERSON \setminus men$

husband ∈ *women* →→ *men*

mother \in *PERSON* \rightarrow dom(*husband*)

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dom(*husband*)



• The usual (+, -, *, \div) plus: mod, $\hat{}$ (power).

• card(set), min(set), max(set)

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